

# Maximum Likelihood Laplacian Correlation Channel Estimation in Layered Wyner-Ziv Coding

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**Abstract**—The distributed source coding theory teaches that efficient compression can be achieved by exploiting the correlation between encoded sources at the decoder only. A critical issue in this coding paradigm is how to accurately estimate the correlation statistics at the decoder. This paper proposes a novel correlation channel estimation method designed for generic layered Wyner-Ziv coding. Driven by real-life applications, i.e., video coding, the correlation noise is assumed to be memoryless zero-mean Laplacian with a block-stationary behavior. Unlike existing methods, the proposed technique derives a novel maximum likelihood estimate of the noise scaling parameter per stationarity block. Adhering to layered coding, the derived estimate is successively refined as more stages are decoded. The Crámer-Rao lower bound (CRLB) for our estimator is derived and the mean squared error performance of the proposed successive refinement algorithm is evaluated through simulations. Apart from achieving an accuracy close to the derived CRLB, the proposed algorithm is shown to yield a Wyner-Ziv rate-distortion performance approaching the ideal but unrealistic scenario where the decoder knows *a priori* the correlation statistics. Moreover, the proposed technique has been integrated into our latest Wyner-Ziv video codec for wireless capsule endoscopy. Experimentation reveals that the proposed method delivers improved compression performance compared to state-of-the-art techniques, while simultaneously reducing the overall decoding complexity.

**Index Terms**—Layered Wyner-Ziv (WZ) coding, correlation channel estimation (CCE), maximum likelihood estimate (MLE), Crámer-Rao lower bound (CRLB), wireless capsule endoscopy.

## I. INTRODUCTION

**D**ISTRIBUTED source coding (DSC) refers to the separate encoding and joint decoding of correlated sources. Slepian and Wolf [1] derived the theoretically achievable rate region for the distributed coding of two discrete sources in a lossless compression scenario. Wyner and Ziv [2] established

the lower rate bound under a distortion constraint when a source is encoded independently but decoded using a correlated signal as side information. In successively refined Wyner-Ziv (WZ) coding [3], [4], a source is encoded in multiple stages and decoded with different distortion and possibly different side information at each stage.

Expressing the statistical dependency between the distributed sources in the form of a virtual channel, one may design Slepian-Wolf (SW) coding based on binning realized by channel codes [5]–[7]. Alternatively, distributed arithmetic codes [8], [9] can also implement SW coding. Practical WZ constructions include nested lattice or trellis codes [10] and SW coded nested [11] or non-nested [6] lattice quantization. In [4], successively refined WZ coding was constructed by nested scalar quantization followed by SW coding of bit-planes with low-density parity-check (LDPC) codes.

DSC finds applications in several domains, including wireless sensor networks [12], biometrics, and tampering localization [13]. Distributed video coding (DVC) [6], [12], [14], in particular, targets lightweight multimedia applications, such as wireless visual sensors and wireless capsule endoscopy [15]. In this context, layered WZ coding [4] provides scalability, a key trait for transmission over heterogeneous wireless networks with varying resources and bandwidth.

A common assumption in DSC schemes, e.g., [4], [10], [11], [16], is that the joint decoding of the sources is carried out with a priori knowledge of the correlation statistics. However, in real-world signals, for instance, obtained by sensors that measure temperature, humidity, or capture sound or visual data, the correlation statistics vary both spatially and temporally. The varying correlation statistics cannot be directly measured since the sources are independently encoded. Hence, accurate correlation channel estimation (CCE) is critical to achieve high joint decoding performance.

### A. Related Work

Various works have addressed CCE for SW coding of binary sources, where the correlation is modeled as a binary symmetric channel (BSC). To estimate the crossover probability of the BSC, residual redundancies in the log-likelihood ratios were exploited in [17], while expectation-maximization (EM) algorithms were proposed in [18], [19]. In [20], a maximum likelihood estimate (MLE) was obtained using the source and the side information syndromes, which served as an initialization for the EM algorithm. Nevertheless, in [17]–[20] the correlation noise was assumed to be stationary within a

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codeword. A CCE approach that can track the varying BSC crossover parameter in a codeword by incorporating particle filters in an LDPC factor graph, was proposed in [21]. Particle filtering was subsequently substituted by an expectation-propagation algorithm, resulting in comparable accuracy and reduced computational overhead [22].

Concerning WZ coding, several CCE methods have been proposed in the context of DVC. The majority of the schemes assume an additive noise channel, where the correlation noise is an independent memoryless zero-mean Laplacian distributed random variable [6], [23]. Several stationarity levels of the correlation noise have been considered. In transform-domain DVC, for instance, the correlation noise can be estimated for every frequency band per frame or even per transform coefficient [23]. To improve the CCE in the transform domain, the use of a pixel-domain CCE and the spatial correlation of the noise signal is proposed in [24]. A similar concept was exploited in the TRACE algorithm [25], where refinement of the noise variance upon decoding of each DCT band was carried out. Exploiting cross-band and inter-bit-plane dependencies [26] for noise frame residue refinement was put forward in [27]. Alternatively, block-based classification of the Laplace scaling parameter using offline training was performed in [28]. An adaptive CCE method based on clustering of DCT blocks, updated at the bit-plane level, was presented in [29]. Nevertheless, these CCE techniques for DVC are application-specific; many of which are even explicitly tailored to particular codec components; e.g., the methods in [23], [24], [27]–[29] are designed for motion-compensated interpolation based side information creation.

An application-generic approach to CCE for WZ coding, the method in [30] exploits particle filtering based belief propagation, enabling adaptation to fine noise stationarity levels. However, it does not support layered coding since all the bit-planes are decoded on a joint factor graph. In addition, the increased code length and the use of particle filtering put a high computational strain on the decoder. Joint bit-plane decoding was also used in [31] as to exploit inter-bit-plane correlation in the WZ decoding and in the CCE refinement.

Yet, only few CCE solutions can address application-generic layered WZ coding. The authors of [32] extended the method in [20] to count for the inter-bit-plane dependencies expressed using a Laplace correlation noise model. In prior work [33], we proposed an efficient bit-plane-per-bit-plane successively refined CCE for the side-information-dependent (SID) correlation channel. In [32], [33], however, the correlation noise was assumed stationary per DCT band of each frame.

## B. Contributions

Targeting at efficient generic layered WZ coding, this paper proposes a novel approach to progressively refined maximum likelihood (ML) estimation of the Laplacian correlation channel. The proposed technique features key novelties with respect to the state-of-the-art.

First, the proposed method considers the correlation noise statistics to vary per block (a.k.a. chunk) of consecutive samples in a codeword and provides accurate CCE even when

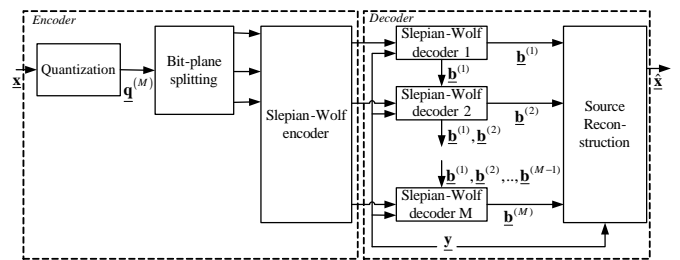


Fig. 1. The layered Wyner-Ziv coding scheme.

small noise stationarity blocks are assumed. This opposes the methods in [32], [33], where the correlation noise was considered stationary within a codeword.

Second, at each WZ decoding stage, we derive a novel estimate of the correlation noise scaling parameter per stationarity block by exploiting the side information and the decoding information from the previous decoded stages. Conversely to our previous work [33], it is proven that the derived estimate constitutes the MLE of the correlation noise parameter. We also derive the Crámer-Rao lower bound (CRLB) on the accuracy of our MLE per decoding stage. We show that, as the number of decoded stages increases, the derived CRLB approaches the oracle-based CRLB, that is, the ideal bound attained when knowing the correlation noise samples.

Third, unlike methods solely designed for DVC, e.g., [23]–[25], [27], [28], the proposed approach is application-generic. When applied to DVC, particularly, the proposed technique is neither confined to a specific side information generation method (unlike [23], [24], [27]) nor does it impose a fixed decoding order of the DCT bands like the approach in [25].

Simulations on synthetic data show that, per decoding stage, the mean squared error (MSE) of the proposed MLE is vastly decreased compared to our previous state-of-the-art estimator from [33] and approaches the derived CRLB. Furthermore, the WZ rate-distortion (RD) performance obtained by using the proposed method is close to the one obtained by assuming perfect knowledge of the noise statistics. When implemented in our state-of-the-art DVC system for wireless capsule endoscopy, the proposed method yields systematic compression gains over the TRACE technique in [25] and our prior SID CCE method in [33] (i.e., rate savings up to 7.13% and 4.22%, respectively). The results also show that the proposed algorithm causes a reduction of the decoding complexity.

Concerning the remainder of the paper, Section II elaborates on layered WZ coding and states the CCE problem. The proposed CCE algorithm is presented in Section III, while bounds on the estimation accuracy are derived in Section IV. Section V expands on the application of the proposed CCE in wireless capsule endoscopy, and Section VI reports our experimental results. Finally, Section VII concludes the work.

## II. PROBLEM SETTING

### A. Layered Wyner-Ziv Coding Architecture

We consider layered WZ coding of a source  $X$  given correlated side information  $Y$  available at the decoder. Contrary

to [4], in the considered code design (see Fig. 1), source clustering is realized by scalar quantization without nesting. The design is motivated by constructions [34] that employ classical quantization and exclusively assign the binning operation to a highly-dimensional channel code. Such code constructions are commonly used in DVC, e.g., [6], [15], [25], [32].

1) *Encoder*: Consider a quantization function  $Q_M(X)$ ,  $Q_M : \mathcal{X} \rightarrow \{1, 2, \dots, 2^M\}$ , that is a mapping from the source to the quantization alphabet. Given the partition  $Q_M$ , the encoder quantizes the incoming source samples  $n$ -tuple  $\mathbf{x} \triangleq x_1, x_2, \dots, x_n$ ,  $x_i \in \mathcal{X}$ , to a quantization index  $n$ -tuple  $\mathbf{q}^{(M)} \triangleq q_1^{(M)}, q_2^{(M)}, \dots, q_n^{(M)}$ ,  $q_i^{(M)} \in \{1, 2, \dots, 2^M\}$ . The quantization index  $n$ -tuple  $\mathbf{q}^{(M)}$  is then divided into a number of  $M$  bit-planes of length  $n$ , i.e.,  $\mathbf{b}^{(1)}, \mathbf{b}^{(2)}, \dots, \mathbf{b}^{(M)}$ , where  $\mathbf{b}^{(m)} \triangleq b_1^{(m)}, b_2^{(m)}, \dots, b_n^{(m)}$  denotes the  $m$ -th bit-plane and  $b_i^{(m)} = \{0, 1\}$  represents the  $i$ -th bit in the  $m$ -th bit-plane.

Using a predefined order—e.g., from the most-significant bit-plane (MSB),  $m = 1$ , to the least significant bit-plane (LSB),  $m = M$ ,—the source bit-planes  $\mathbf{b}^{(1)}, \mathbf{b}^{(2)}, \dots, \mathbf{b}^{(M)}$  are sequentially passed to a SW encoder that can ideally achieve the compression rate  $H(Q_M(X)|Y)$ . In this work, structured SW binning is implemented by the rate-adaptive LDPC accumulate (LDPCA) codes of [5].

2) *Decoder*: The decoder exploits the correlation statistics to sequentially decode the source bit-planes  $\mathbf{b}^{(1)}, \mathbf{b}^{(2)}, \dots, \mathbf{b}^{(M)}$  with the aid of the side information samples  $n$ -tuple  $\mathbf{y}$ . In particular, for every source bit-plane  $\mathbf{b}^{(m)}$ ,  $m = \{1, 2, \dots, M\}$ , the correlation statistics between the source  $n$ -tuple  $\mathbf{x}$  and the side information  $n$ -tuple  $\mathbf{y}$  are converted to soft-input estimates, namely, log-likelihood ratios (LLRs). These LLRs, which provide *a priori* information about the probability of each bit to be  $b_i^{(m)} = 0$  or  $b_i^{(m)} = 1$ , are passed to the soft channel decoder. Then, an iterative soft-decoding algorithm (e.g., the sum-product algorithm [35]) is executed to decode the source bit-planes with an error probability close to zero.

For optimal layered WZ coding, during the formulation of the LLRs, information given by the side information and the already decoded source bit-planes is taken into account. Specifically, let  $b_i^{(m)}$ ,  $m \in \{1, 2, \dots, M\}$ ,  $i \in \{1, 2, \dots, n\}$ , be the bit under consideration,  $y_i$  be the corresponding side information sample value and  $b_i^{(1)}, b_i^{(2)}, \dots, b_i^{(m-1)}$  be the already decoded bits in the previous  $m-1$  bit-planes. Then, the input LLR is given by  $\log \frac{P(b_i^{(m)}=0|y_i, b_i^{(1)}, b_i^{(2)}, \dots, b_i^{(m-1)})}{P(b_i^{(m)}=1|y_i, b_i^{(1)}, b_i^{(2)}, \dots, b_i^{(m-1)})}$ , where the numerator and the denominator are calculated by integrating the conditional probability density function (pdf) of the correlation channel, i.e.,  $f_{X|Y}(x|y)$ , over the quantization bin defined by the already decoded bits  $b_i^{(1)}, b_i^{(2)}, \dots, b_i^{(m-1)}$  and by setting  $b_i^{(m)} = 0$  or  $b_i^{(m)} = 1$ , respectively. Because of this chained approach to construct the LLRs, the total rate conforms to the chain rule of entropies [36]. This means that, ideally, there is no performance loss between non-layered and layered WZ coding [4].

Once all the source bit-planes are decoded, they are combined to form the decoded quantization indices of the source. These indices are then used to reconstruct the source with the help of the side information and the correlation statistics. In

particular, if the mean squared error (MSE) distortion metric is employed, the optimal reconstruction of a source sample  $x_i$  is obtained as the centroid of the random variable  $X$  given the corresponding side information sample  $y_i$  and the decoded quantization index  $q_i^{(M)}$ . Namely,  $\hat{x}_i = \mathbb{E} \left[ x_i | y_i, q_i^{(M)} \right]$ , where  $\mathbb{E}[\bullet|\bullet]$  is the conditional expectation operator. Closed-form expressions for the reconstruction function can be found in [4] and [37] for a Gaussian and a Laplacian pdf, respectively.

3) *Rate Control*: Early DSC designs [4], [10], [11], assume the correlation statistics to be stationary and perfectly known *a priori* at the encoder and decoder. Under this assumption, the encoder and decoder agree on an efficient code rate driven by the SW rate. Yet, in practice, the correlation channel exhibits a non-stationary behavior, which neither can be known in advance nor can it be measured in practice. Therefore, a rate-adaptive decoder-driver rate control approach using a feedback channel is a good solution for many applications [6], [15], [32].

## B. The Correlation Channel Estimation Problem

We assume that the correlation between the source  $X$  and the side information  $Y$  is expressed by a memoryless channel,  $X = Y + Z$ , where the correlation noise  $Z$  is zero-mean Laplacian, i.e.,  $Z \sim \text{Laplace}(0, \lambda)$ , and independent of the side information, which is arbitrarily distributed. The assumption of independent correlation noise opposes our previous work [33], where the noise was dependent on the side information. Our motivation stems from the fact that most existing DVC systems employ the former model to express the correlation channel [6], [14], [23]–[32]. The correlation channel pdf is then given by

$$f_{X|Y}(x|y) = \frac{\lambda}{2} \exp(-\lambda|x-y|), \quad (1)$$

where  $\lambda = \frac{\sqrt{2}}{\sigma_Z}$  is the scaling parameter of the distribution and  $\sigma_Z^2 = \mathbb{E}\{Z^2\}$ . We assume that the correlation noise  $Z$  is *block-stationary*; namely,  $\lambda$  varies per noise stationarity block (chunk), defined by  $K$  consecutive samples in the noise  $n$ -tuple  $\mathbf{z}$ , where  $n \bmod K = 0$ . To express this formally, we adhere to the following notation. An  $n$ -tuple  $\mathbf{z}$  is written as  $\mathbf{z} \triangleq \mathbf{z}_0 \mathbf{z}_1 \mathbf{z}_2 \dots \mathbf{z}_{L-1}$ , where  $\mathbf{z}_\ell$  denotes a block of the  $K$  succeeding noise samples  $z_{\ell \times K+i}$ , where  $i = \{1, 2, \dots, K\}$  counts the samples in the block,  $\ell \in \{0, 1, \dots, L-1\}$  indexes the blocks, and  $L \triangleq \frac{n}{K}$ . The set  $\{\lambda_\ell\}_{\ell=0}^{L-1}$  contains the scaling parameters that control the noise distribution over the entire  $n$ -tuple  $\mathbf{z}$ .

Under the common but unrealistic assumption that the correlation noise samples  $\mathbf{z}$  are known *a priori* to the decoder, an offline MLE for the scaling parameter per stationarity block  $\mathbf{z}_\ell$  is given by

$$\hat{\lambda}_\ell^{\text{off}} = K \left( \sum_{i=1}^K |z_{\ell \times K+i}| \right)^{-1}, \quad \ell = 0, 1, \dots, L-1. \quad (2)$$

We remark that  $\hat{\lambda}_\ell^{\text{off}}$  in (2) is a minimum variance unbiased (MVU) estimator [38].

In practical DSC designs, however, the correlation noise cannot be directly measured as the source and the side information are located at the encoder and decoder, respectively.

In the following, we derive a novel online MLE for  $\lambda_\ell$  per stationary block at the decoder, which can be successively refined as source bit-planes are decoded. We also evaluate the asymptotic efficiency of the proposed MLE using the CRLB.

We remark that a block-stationary noise model with a constant number of samples per stationarity block is assumed for simplicity. The focus of the paper is on the derivation and the evaluation of the performance of the MLE  $\lambda_\ell$ . Devising mechanisms to adapt to complex noise variations through arbitrary shaped blocks is outside the scope of this paper.

### III. MAXIMUM LIKELIHOOD CORRELATION ESTIMATION

#### A. Introduction to the Basic Concept

In the layered WZ coding scheme, discussed in Section II-A, the source  $n$ -tuple  $\mathbf{x}$  is quantized with a  $2^M$ -stage quantization partition  $Q_M$  and successively encoded into  $M$  bit-planes. This progressive coding scheme is equivalent to an *embedded* partition  $Q_1 Q_2 \dots Q_m \dots Q_M$ , where the  $m$ -th stage partition  $Q_m$  has  $2^m$  quantization intervals corresponding to indices  $q_i^{(m)} \in \{1, 2, \dots, 2^m\}$ , formed by the bits  $b_i^{(1)} b_i^{(2)} \dots b_i^{(m)}$ . We exploit this progressive coding approach at the decoder to design a successively refined MLE of the *block-stationary* correlation channel parameter  $\lambda_\ell$ .

The proposed progressive ML CCE scheme is described as follows: Starting from initial correlation channel parameter estimates  $\{\hat{\lambda}_\ell^{(0)}\}_{\ell=0}^{L-1}$ , the SW decoder decodes the MSB  $\mathbf{b}^{(1)}$  of the source samples. This initial estimate may for instance be obtained from a previously decoded data  $n$ -tuple, or it can be the expected value of the parameter derived from offline observations.

By decoding  $\mathbf{b}^{(1)}$ , the decoder obtains access to a coarsely quantized version of the source samples, denoted by the  $n$ -tuple  $\mathbf{q}^{(1)}$ , which corresponds to partition  $Q_1$ . The decoder can exploit the available quantized source  $n$ -tuple  $\mathbf{q}^{(1)}$  and the side information  $n$ -tuple  $\mathbf{y}$  to perform CCE for the next bit-plane. Analogous to the decomposition of the noise samples into stationarity blocks, let us write  $\mathbf{q}^{(1)} \triangleq \mathbf{q}_0^{(1)} \mathbf{q}_1^{(1)} \mathbf{q}_2^{(1)} \dots \mathbf{q}_{L-1}^{(1)}$  and  $\mathbf{y} \triangleq \mathbf{y}_0 \mathbf{y}_1 \mathbf{y}_2 \dots \mathbf{y}_{L-1}$ . Namely, per stationarity block  $\ell$ , the decoder derives an MLE  $\hat{\lambda}_\ell^{(1)}$  using the samples  $q_{\ell \times K+i}^{(1)}$  and  $y_{\ell \times K+i}$ ,  $i \in \{1, 2, \dots, K\}$ , in the  $K$ -tuples  $\mathbf{q}_\ell^{(1)}$  and  $\mathbf{y}_\ell$ , respectively. The proposed derivation of the MLE is detailed in Section III-B. The updated parameters  $\{\hat{\lambda}_\ell^{(1)}\}_{\ell=0}^{L-1}$  are then used to derive the soft-input information to SW decode the next source bit-plane  $\mathbf{b}^{(2)}$ . The decoder then combines the available  $\mathbf{b}^{(1)}$  and  $\mathbf{b}^{(2)}$  to form a more accurate version, i.e.,  $\mathbf{b}^{(2)}$ , of the quantized source samples (corresponding to partition  $Q_2$ ), which is in turn used to obtain a refined MLE  $\hat{\lambda}_\ell^{(2)}$  per stationarity block  $\ell$ .

In a recursive manner, after SW decoding of the  $\mathbf{b}^{(m)}$  bit-plane ( $1 \leq m \leq M$ ), the proposed algorithm forms a quantized version  $\mathbf{q}^{(m)}$  of the source samples based on the available bit-planes  $\mathbf{b}^{(1)} \mathbf{b}^{(2)} \dots \mathbf{b}^{(m)}$ . Using  $\mathbf{q}_\ell^{(m)}$  and  $\mathbf{y}_\ell$  the decoder obtains an MLE  $\hat{\lambda}_\ell^{(m)}$  for the  $\ell$ -th stationarity block. We highlight that when all bit-planes have been decoded, a.k.a.,  $m = M$ , the proposed MLE algorithm is carried out

again so as to obtain the channel parameters  $\{\hat{\lambda}_\ell^{(M)}\}_{\ell=0}^{L-1}$ , which are used to perform optimal minimum MSE (MMSE) reconstruction. Since per additional decoded bit-plane  $\mathbf{b}^{(m)}$  the available approximation  $\mathbf{q}^{(m)}$  of the source samples at the decoder becomes more accurate, the obtained MLE  $\{\hat{\lambda}_\ell^{(m)}\}_{\ell=0}^{L-1}$  are becoming more accurate as well, thereby leading to a progressively improved CCE.

#### B. Maximum Likelihood Estimation Formulations

We now concentrate on the proposed derivation of the MLE  $\lambda_\ell^{(m)}$  from the samples in the  $K$ -tuples  $\mathbf{q}_\ell^{(m)}$  and  $\mathbf{y}_\ell$ . Since we focus on the samples in a stationarity block and in order to maintain the simplicity of the notations, without loss of generality, we hereby drop the subscript  $\ell$  that indexes the stationarity block. Consider the following definition.

*Definition 1 (m-th Stage Estimator):* Let  $\mathbf{q}^{(m)} \triangleq q_1^{(m)}, q_2^{(m)}, \dots, q_K^{(m)}$ ,  $\mathbf{q}^{(m)} \in \{1, 2, \dots, 2^m\}^K$  be a  $K$ -tuple of quantization indices of the WZ source samples, formed by  $m$  SW decoded bit-planes,  $1 \leq m \leq M$ . Also, let  $\mathbf{y} \triangleq y_1, y_2, \dots, y_K$  be a  $K$ -tuple of side information samples with  $\mathbf{y} \in \mathcal{Y}^K$ . An  $m$ -th stage estimator  $\hat{\lambda}^{(m)}$  of the Laplace correlation channel in (1) is defined as a function  $\varphi: \{1, 2, \dots, 2^m\}^K \times \mathcal{Y}^K \rightarrow \Lambda$ , where  $\Lambda = (0, +\infty)$  is the parameter set for  $\hat{\lambda}^{(m)}$  and  $K$  is the samples' size.

We consider the  $m$ -th stage *conditional likelihood function*  $\mathcal{L}(\lambda) = p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$ , where  $p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$  is the conditional probability mass function (pmf) of the quantization indices  $K$ -tuple given the side information samples'  $K$ -tuple, parameterized by the parameter  $\lambda$  of the correlation channel pdf. Since the correlation channel is assumed memoryless, we have

$$\mathcal{L}(\lambda) = p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = \prod_{i=1}^K p(q_i^{(m)}|y_i; \lambda), \quad (3)$$

where  $p(q_i^{(m)}|y_i; \lambda)$ ,  $\forall i \in \{1, 2, \dots, K\}$ , is the conditional pmf for the quantization index  $q_i^{(m)}$  of a source sample  $x_i$  given its corresponding side information sample  $y_i$ . The conditional pmf  $p(q_i^{(m)}|y_i; \lambda)$  is defined by the integration of the correlation channel pdf  $f_{X|Y}(x|y)$  over the quantization bin defined by  $q_i^{(m)}$ , that is,

$$\begin{aligned} p(q_i^{(m)}|y_i; \lambda) &= \int_{u_i^{(m)}}^{v_i^{(m)}} f_{X|Y}(x|y_i) dx \\ &= \int_{u_i^{(m)}}^{v_i^{(m)}} \frac{\lambda}{2} \exp(-\lambda|x - y_i|) dx, \end{aligned} \quad (4)$$

where  $u_i^{(m)}$  and  $v_i^{(m)}$  are respectively the lower and upper bound of the quantization bin indexed by  $q_i^{(m)}$  in the quantization partition  $Q_m$ . Depending on the relative position of  $y_i$  with respect to the bounds  $[u_i^{(m)}, v_i^{(m)})$  of the bin indexed by

$q_i^{(m)}$ , the probability  $p(q_i^{(m)}|y_i; \lambda)$  can be computed as

$$p(q_i^{(m)}|y_i; \lambda) = \begin{cases} \frac{1}{2}e^{-\lambda(u_i^{(m)}-y_i)} - \frac{1}{2}e^{-\lambda(v_i^{(m)}-y_i)}, & y_i < u_i^{(m)} \\ \frac{1}{2}e^{\lambda(v_i^{(m)}-y_i)} - \frac{1}{2}e^{\lambda(u_i^{(m)}-y_i)}, & y_i \geq v_i^{(m)} \\ 1 - \frac{1}{2}e^{\lambda(u_i^{(m)}-y_i)} - \frac{1}{2}e^{-\lambda(v_i^{(m)}-y_i)}, & u_i^{(m)} \leq y_i < v_i^{(m)} \end{cases} \quad (5)$$

We wish to find an  $m$ -th stage MLE  $\hat{\lambda}^{(m)}$ , such that

$$\hat{\lambda}^{(m)} = \arg \max_{\lambda} \mathcal{L}(\lambda). \quad (6)$$

To ease the calculations, we derive  $\hat{\lambda}^{(m)}$  by equivalently maximizing a monotone transformation of (3), that is

$$\ln \mathcal{L}(\lambda) = \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = \sum_{i=1}^K \ln p(q_i^{(m)}|y_i; \lambda). \quad (7)$$

Then, the  $m$ -th stage MLE  $\hat{\lambda}^{(m)}$  that achieves the maximum in (6), can be found as the solution of the following  $m$ -th stage *conditional log-likelihood equation*

$$\frac{\partial}{\partial \lambda} \ln \mathcal{L}(\lambda) = 0 \Rightarrow \sum_{i=1}^K \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = 0, \quad (8)$$

where the partial derivative  $\frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda)$  is derived in Appendix A for each of the three forms of (5), corresponding to the cases  $y_i < u_i^{(m)}$ ,  $y_i \geq v_i^{(m)}$  and  $u_i^{(m)} \leq y_i < v_i^{(m)}$ . Concerning the solution of (8), the following holds.

**Theorem 1 (Existence and Uniqueness):** There is a unique finite  $m$ -th stage MLE  $\hat{\lambda}^{(m)}$  of the Laplace correlation channel pdf  $f_{X|Y}(x|y)$  for  $K$  samples, under the condition that there exists at least one side information sample  $y_i$  in  $\mathbf{y}$  that is not between or on the bounds of the quantization bin defined by its corresponding  $m$ -th stage quantization index  $q_i^{(m)}$ .

*Proof:* The proof of Theorem 1 is given in Appendix B preceded by a preparatory lemma. ■

**Corollary 1:** If every side information sample  $y_i$  in  $\mathbf{y}$  is located between or on the bounds of the quantization bin defined by its corresponding  $m$ -th stage quantization index  $q_i^{(m)}$ , namely, if  $u_i^{(m)} \leq y_i \leq v_i^{(m)}$ ,  $\forall i \in \{1, 2, \dots, K\}$ , then the  $m$ -th stage MLE is given by  $\hat{\lambda}^{(m)} = +\infty$ .

*Proof:* Intuitively, when  $u_i^{(m)} \leq y_i \leq v_i^{(m)} \forall y_i$  in  $\mathbf{y}$ , then the likelihood is maximized at  $\sigma_Z = 0$  or  $\lambda = +\infty$ . Formally, from the proof of Theorem 1, if  $u_i^{(m)} \leq y_i \leq v_i^{(m)}$ ,  $\forall i \in \{1, 2, \dots, K\}$  then  $\lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = 0$ . Combining this result with Lemma 1 in Appendix B ends the proof. ■

The following two propositions describe the conditions under which an analytical derivation of the finite MLE, defined in Theorem 1, is feasible.

**Proposition 1 (Analytical Solutions):** The MLE can be found analytically if the following sufficient conditions are satisfied: (i) the quantizer has uniform intervals and (ii)  $y_i \notin [u_i^{(m)}, v_i^{(m)})$ ,  $\forall i = \{1, 2, \dots, K\}$ .

*Proof:* The proof is sketched in Appendix C. ■

**Proposition 2 (High Rate Assumptions):** Under high rate assumptions, the MLE is approximated by  $\hat{\lambda}^{(m)} = K \left( \sum_{i=1}^K |\xi_i^{(m)} - y_i| \right)^{-1}$ , where  $\xi_i^{(m)} = \frac{u_i^{(m)} + v_i^{(m)}}{2}$ .

*Proof:* Let  $\Delta_i^{(m)} = v_i^{(m)} - u_i^{(m)}$  denote the quantization interval at the  $m$ -th decoding stage. When  $\Delta_i^{(m)} \rightarrow 0$  the pdf  $f_{X|Y}(x|y)$  is assumed to be constant in the quantization interval and thus, the pmf in (4) can be approximated as  $p(q_i^{(m)}|y_i; \lambda) \simeq \frac{\lambda}{2} e^{-\lambda|\xi_i^{(m)}-y_i|} \Delta_i^{(m)}$ . Inserting the latter in (8) and solving for  $\lambda$  ends the proof. ■

**Remark 1:** In the general case, the conditional log-likelihood equation  $\frac{\partial}{\partial \lambda} \ln \mathcal{L}(\lambda) = 0$  is an exponential equation on  $\lambda$  and cannot be solved analytically. Hence, we hinge on an numerical solution. A grid search for the  $\hat{\lambda}^{(m)}$  that attains the maximum in (6) is not computationally feasible, since the parameter set for  $\lambda$  is continuous and not confined to a finite interval, i.e.,  $\Lambda = (0, +\infty)$ . To derive the MLE, we implement a numerical root finding algorithm that combines bracketing, bisection, secant, and inverse quadratic interpolation methods [39]. After a maximum number of iterations  $J_{MAX}$ , if the algorithm cannot find the root of equation (8), with a given accuracy  $\epsilon_{acc}$ , then we set  $\hat{\lambda}^{(m)} = \hat{\lambda}^{(m-1)}$ . Namely, the correlation channel parameter estimate from the previous decoding level is retained. The same holds for the case that falls into Corollary 1. It is noteworthy that as more bit-planes are decoded the quantization intervals in the  $m$ -th stage partition  $Q_m$  in the decoder become finer and in turn, the probability of having a case corresponding to Corollary 1 drops. This effect is in accordance with the progressively refined nature of the proposed estimator.

The proposed ML CCE algorithm, as applied for the  $\ell$ -th stationarity block, is summarized in Algorithm 1.

---

#### Algorithm 1 ML Estimation at Stationary Block Level

---

```

1: function MLE( $\mathbf{q}_\ell^{(m)}, \mathbf{y}_\ell$ )
2:   if  $u_{\ell \times K+i}^{(m)} \leq y_{\ell \times K+i} \leq v_{\ell \times K+i}^{(m)}, \forall i \in \{1, 2, \dots, K\}$ 
3:     then
4:       Set  $\hat{\lambda}_\ell^{(m)} = \hat{\lambda}_\ell^{(m-1)}$ 
5:     else
6:       if  $y_{\ell \times K+i} \notin [u_{\ell \times K+i}^{(m)}, v_{\ell \times K+i}^{(m)}) \wedge$ 
7:          $\Delta_{\ell \times K+i}^{(m)} = \Delta^{(m)}, \forall i \in \{1, 2, \dots, K\}$ 
8:         then
9:            $\hat{\lambda}_\ell^{(m)}$  is analytically found by (18) or (19).
10:        else
11:          Find  $\hat{\lambda}_\ell^{(m)}$  by numerically solving the equation
12:           $\frac{\partial}{\partial \lambda} \ln \mathcal{L}(\lambda) = 0$  defined in (8).
13:        end if
14:      end if
15:    end if
16:  end function
    
```

---

**Remark 2:** For simplicity, the above presentation of the ML estimation formulations consider identical side information samples  $\mathbf{y}$  at every refinement level  $m \in \{1, 2, \dots, M\}$ . Yet, since the proposed ML CCE algorithm is applied per refinement stage  $m$ , it can be directly generalized to the case where the side information samples  $\mathbf{y}^{(m)}$  are refined per level.

*Remark 3:* The proposed progressively refined ML CCE algorithm is applicable irrespective of the side information's distribution and the quantization function used at the layered WZ encoder. The MLE is derived based on the side information samples and the bounds  $[u_i^{(m)}, v_i^{(m)})$  of the quantization bin corresponding to each quantized source sample  $q_i^{(m)}$ .

We conjecture that the proposed technique can be extended to estimate the standard deviation of a zero-mean Gaussian correlation noise as well. The proof and validation of this conjecture is left as topic for further research.

#### IV. ASYMPTOTIC EFFICIENCY BOUNDS

The  $m$ -th stage MLE, which results from solving (8), is biased since (8) is a non-linear equation on  $\lambda$ . To study the asymptotic efficiency of the proposed MLE, we derive the  $m$ -th stage CRLB<sup>( $m$ )</sup> for the problem. We observe that the domain in which the  $m$ -th stage conditional likelihood function in (3) is nonzero does not depend on  $\lambda$ . Hence, the regularity condition  $\mathbb{E} \left[ \frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) | Y \right] = 0$  is satisfied [38]. As a consequence, the CRLB<sup>( $m$ )</sup> is given by<sup>1</sup>

$$\text{CRLB}^{(m)} = -\mathbb{E} \left[ \frac{\partial^2}{\partial \lambda^2} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) | Y \right]^{-1}, \quad (9)$$

where  $\mathbb{E}[\bullet|Y]$  is the conditional expectation with respect to the conditional pmf  $p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$ . The average curvature of the conditional log-likelihood function can be computed as

$$\begin{aligned} & \mathbb{E} \left[ \frac{\partial^2}{\partial \lambda^2} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) | Y \right] \\ &= \sum_{\mathbf{q}^{(m)}} \frac{\partial^2}{\partial \lambda^2} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) \\ &= \sum_{q_1^{(m)}=1}^{2^m} \dots \sum_{q_K^{(m)}=1}^{2^m} \sum_{i=1}^K \frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda) \prod_{i=1}^K p(q_i^{(m)}|y_i; \lambda) \\ &= \sum_{i=1}^K \sum_{q_i^{(m)}=1}^{2^m} \frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda) p(q_i^{(m)}|y_i; \lambda), \quad (10) \end{aligned}$$

where the forms of  $p(q_i^{(m)}|y_i; \lambda)$  and  $\frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda)$  are given by (5) and (12a)-(12c), respectively. Lemma 1 in Appendix B proves that  $\frac{\partial^2}{\partial \lambda^2} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) < 0$ , thereby verifying that CRLB<sup>( $m$ )</sup>  $> 0$ . For the analytical forms of the MLE, obtained under the conditions described in Proposition 1, the  $m$ -th stage CRLB<sup>( $m$ )</sup>'s are given in Appendix C.

It is worth noticing that, as the number of decoded WZ stages increases, the CRLB<sup>( $m$ )</sup> asymptotically (i.e.,  $\Delta_i^{(m)} \rightarrow 0$ ) approaches the oracle (or offline) CRLB. The latter, which is defined by  $\frac{\lambda^2}{K}$ , refers to the unrealistic case where the decoder has access to the noise samples. This convergence to the oracle CRLB can be readily verified by replacing the approximated high rate pmf (given in the proof of Proposition 2) into (10).

<sup>1</sup>For the sake of simplicity, we again drop the subscript  $\ell$  that indexes the stationarity block size.

#### V. APPLICATION TO WYNER-ZIV CODING OF ENDOSCOPIC VIDEO

The practical merit of the proposed MLE is shown using our endoscopic DVC (EDVC) [15] system. EDVC delivers state-of-the-art DVC performance, outperforming several relevant up-to-date codecs, including [14], [33], [40], [41].

##### A. Codec Architecture Overview

In EDVC [15], the key frames are intra-coded with JPEG or H.264/AVC Intra, whereas the WZ frames are encoded in a hash and a WZ stage. The hash is a reduced resolution version of each WZ frame coded at a low quality with JPEG or H.264/AVC Intra. Each WZ frame undergoes the discrete cosine transform (DCT) and then the coefficients of each band  $\beta$  are encoded following the codec in Section II-A. Namely, the quantized coefficients per band  $\mathbf{q}^{(M)}(\beta)$  are turned into bit-planes  $\mathbf{b}^{(M)}(\beta)$  that are passed to an LDPCA encoder [5].

In the decoder, the decoded key frames are used together with the decoded and up-scaled hash to produce a motion-compensated prediction of the WZ frame (see [15] for details). Transforming the prediction frame yields side information. In [15], correlation estimation was performed using our SID CCE algorithm from [33] that assumes band-level noise stationarity per frame. In this work, however, we implement our novel ML CCE, which is applied for each stationarity block per band of each frame. The obtained statistics are used to decode the bit-planes of the WZ bands and perform reconstruction. Inverse integer DCT is finally carried out to return to the pixel domain.

##### B. Proposed ML Correlation Estimation Technique in EDVC

The correlation between the DCT coefficients of a WZ frame and its side information is considered Laplacian with a scaling parameter varying per group of neighboring coefficients in each band per frame. For consistency, following the notations from Section II, the size of a band of a frame is denoted by  $n = V \times H$ , where  $V$ ,  $H$  represent the number of DCT blocks in the vertical and horizontal direction of a transformed frame, respectively. Coefficients of the same band, say  $\beta$ , that belong to  $K = K_1 \times K_2$  neighboring DCT blocks are clustered into correlation noise stationarity blocks, where  $K_1, K_2$  are divisors of  $V, H$ , respectively. By definition, there are  $L = \frac{n}{K_1 \times K_2}$  stationarity blocks of coefficients per band  $\beta$  of a frame; namely, for a side information frame band we can write  $\mathbf{y}(\beta) \triangleq \mathbf{y}_0(\beta)\mathbf{y}_1(\beta)\mathbf{y}_2(\beta)\dots\mathbf{y}_{L-1}(\beta)$ . The  $K = K_1 \times K_2$  coefficients in each block  $\mathbf{y}_\ell(\beta)$  share the same correlation noise parameter  $\lambda_\ell(\beta)$ . The set  $\{\lambda_\ell(\beta)\}_{\ell=0}^{L-1}$  refers to the scaling parameters of the noise over the band  $\beta$  of a frame.

The proposed ML CCE technique is executed per band  $\beta$  of each WZ frame. The scaling parameters  $\{\hat{\lambda}_\ell^{(0)}(\beta)\}_{\ell=0}^{L-1}$  needed to initiate the algorithm are copied from the corresponding band of the previously decoded WZ frame<sup>2</sup>. By decoding each

<sup>2</sup>For the first WZ frame in the sequence a predefined set of scaling parameters is used to decode the first bit-plane per band. Specifically, based on offline measurements, one parameter per band is determined as the average value from all the frames in several test sequences.

bit-plane  $\mathbf{b}^{(m)}(\beta)$  in the band, the decoder learns a recursively updated quantized version  $\mathbf{q}^{(m)}(\beta)$  of the WZ coefficients. Based on the  $K$  coefficients in each  $\mathbf{q}_\ell^{(m)}(\beta)$  and  $\mathbf{y}_\ell(\beta)$ , an updated MLE  $\hat{\lambda}_\ell^{(m)}(\beta)$  is derived, as detailed in Section III-B.

The stationarity block size  $K$ , which is regulated by the parameters  $K_1$  and  $K_2$ , affects the accuracy of the algorithm. Low  $K_1$  and  $K_2$  values (correspondingly, low  $K$  values) provide adjustment to the regional variation of the channel statistics but, on the other hand, challenge accurate statistical inference due to narrow statistical support. In particular, a low value of  $K$  increases the probability of the special case described in Corollary 1, i.e.,  $\hat{\lambda}_\ell^{(m)}(\beta) = +\infty$ . For this reason, the algorithm also estimates a single  $\hat{\lambda}^{(m)}(\beta)$  per band of each frame, using all the band samples in  $\mathbf{y}(\beta)$  and  $\mathbf{q}^{(m)}(\beta)$ . This estimate is used in case an MLE for a stationarity block could not be derived.

## VI. EXPERIMENTAL RESULTS

### A. Evaluation on Synthetic Data

We implement the layered WZ coding system detailed in Section II-A. By definition (see Section II-B), the source is formed as  $X = Y + Z$ , where the side information  $Y$  is considered either Gaussian or Laplacian with a zero mean and  $\sigma_Y = 50$ . The correlation noise samples are drawn from an independent zero-mean Laplacian distribution. Per stationarity block  $\ell = 0, 1, \dots, L - 1$  in a codeword, the noise scaling parameter  $\lambda_\ell = \frac{\sqrt{2}}{\sigma_{Z_\ell}}$  is randomly selected from the interval  $\lambda_\ell \in [0.2020, 2.8284]$  (corresponding to  $\sigma_{Z_\ell} \in [0.5, 6]$ ) with uniform probability. The source sample (codeword) length is set to  $n = 1584$  or  $6336$ . Source quantization is realized by a uniform scalar quantizer with  $2^8$  levels, corresponding to  $M = 8$  bit-planes. To SW decode the first bit-plane of the encoded source  $n$ -tuple and thus, to initiate the proposed MLE algorithm, a fixed scaling parameter value, i.e.,  $\left\{ \hat{\lambda}_\ell^{(0)} \right\}_{\ell=0}^{L-1} = 1.5152$ , is used. This initial value constitutes the mean of the interval from which  $\lambda_\ell$  is drawn. Concerning the configuration of the numerical root-finding algorithm, we set the initial value to 0.01, the maximum number of iterations to  $J_{MAX} = 1000$ , and the accuracy to  $\epsilon_{acc} = 10^{-10}$ . Average results over 400 independent trials are presented.

We begin by evaluating the MSE performance of the proposed MLE per decoded stage (represented by the number  $m$  of SW decoded bit-planes) with respect to the derived  $m$ -th stage CRLB $^{(m)}$  and the oracle CRLB. The MSE performance obtained with the initial estimator (i.e., using the fixed parameter value  $\left\{ \hat{\lambda}_\ell^{(0)} \right\}_{\ell=0}^{L-1} = 1.5152$  to decode all the bit-planes) and our previous SII estimator<sup>3</sup> from [33] are also added in the

<sup>3</sup>The SID estimator in [33] is designed for the SID correlation channel, where the distribution of the correlation noise  $Z = X - Y$  is zero-mean Laplacian with a scaling parameter  $\lambda(y)$  that varies depending on the realization  $y$  of the side information. For practicality [33], a finite number of scaling parameters  $\lambda(\tilde{y}_i)$  are estimated by assigning side information samples to finite indices  $\tilde{y}_i$  via quantization. In this specific experiment, the quantizer is uniform with  $2^7$  levels. According to [33, Lemma 2] one may derive the scaling parameter estimate for the SII channel from the estimates derived for the SID channel as  $\hat{\lambda}_{SII} = \left[ \sum_{\tilde{y}_i} p_{\tilde{Y}}(\tilde{y}_i) \frac{1}{\hat{\lambda}^2(\tilde{y}_i)} \right]^{-1/2}$ , where  $p_{\tilde{Y}}(\tilde{y}_i)$  is the pmf of the quantized side information.

comparison. Four experimental setups are considered, corresponding to a Gaussian or Laplacian side information<sup>4</sup> and to a source length of  $n = 1584$  or  $6336$ . In this experimental setup, we set  $n = K$  (i.e.,  $L = 1$ ). The results in Fig. 2 indicate that the proposed MLE systematically yields a significantly lower MSE performance in comparison to the naive initial estimator as well as our previous successively refined SII estimator from [33]. The vast MSE decrease compared to our prior estimator is due to the fact the proposed estimate is the MLE of the noise scaling parameter. In fact, the MSE of our MLE closely follows the CRLB $^{(m)}$  per decoding stage  $m$ , irrespective of the distribution of the side information or the source length. Hence, the proposed MLE is asymptotically efficient given the decoded information at stage  $m$ . It is important to observe that, as  $m$  increases, the CRLB $^{(m)}$  and the MSE of the proposed MLE progressively converge to the oracle CRLB. This behavior corroborates the successively refined nature of the proposed estimator.

Subsequently, we compare the RD performance of layered WZ coding using the proposed MLE algorithm against the performance obtained when assuming perfect knowledge of the noise scaling parameters at the decoder (offline/oracle estimation). The RD performance when the initial fixed scaling parameter is used to decode all the bit-planes is also assessed. In this experimental setup, the source length is set to 1584 samples, the side information is Laplacian and various stationarity block sizes, that is,  $K = 198, 132, 99$ , or  $44$ <sup>5</sup> are considered. Fig. 3 depicts the distortion—expressed in terms of the signal-to-noise ratio  $\text{SNR} = 10 \log(\sigma_X^2 / (\frac{1}{n} \sum (x_i - \hat{x}_i)^2))$ —versus the rate required for the compression of bit-planes  $m = 2$  to  $8$ . The results show that, in case of large stationarity block sizes [see Figs. 3(a)-(b)], the proposed MLE achieves an RD performance almost identical to offline estimation irrespective of the RD point. In case of smaller stationarity block sizes [see Fig. 3(c)-(d)], a performance loss compared to offline estimation is observed due to the reduced statistical support for accurate derivation of the MLE. Compared to decoding with the fixed initial parameter, the proposed MLE systematically yields higher RD performance. In effect, the obtained improvements are mounting with the rate because the proposed MLE is successively refined as more information is decoded.

### B. Evaluation on Endoscopic Video Data

To assess the impact of the proposed algorithm on the performance of our state-of-the-art EDVC [15] system, we carry out experiments on three conventional endoscopic and five capsule<sup>6</sup> endoscopic test video sequences obtained from clinical examinations performed at the Gastroenterology Clinic of the Universitair Ziekenhuis Brussels. The capsule endoscopic sequences, which visualize diverse areas of the gastrointestinal track of 2 patients, contain 150 frames each, acquired at a

<sup>4</sup>As mentioned in Remark 3, the algorithm is applied unaltered irrespective of the distribution of the side information.

<sup>5</sup>The scaling parameter of the correlation noise is constant for the  $K$  samples within each stationarity block.

<sup>6</sup>The capsule endoscopic video data were obtained with the PillCam SB2 [42] from Given Imaging.

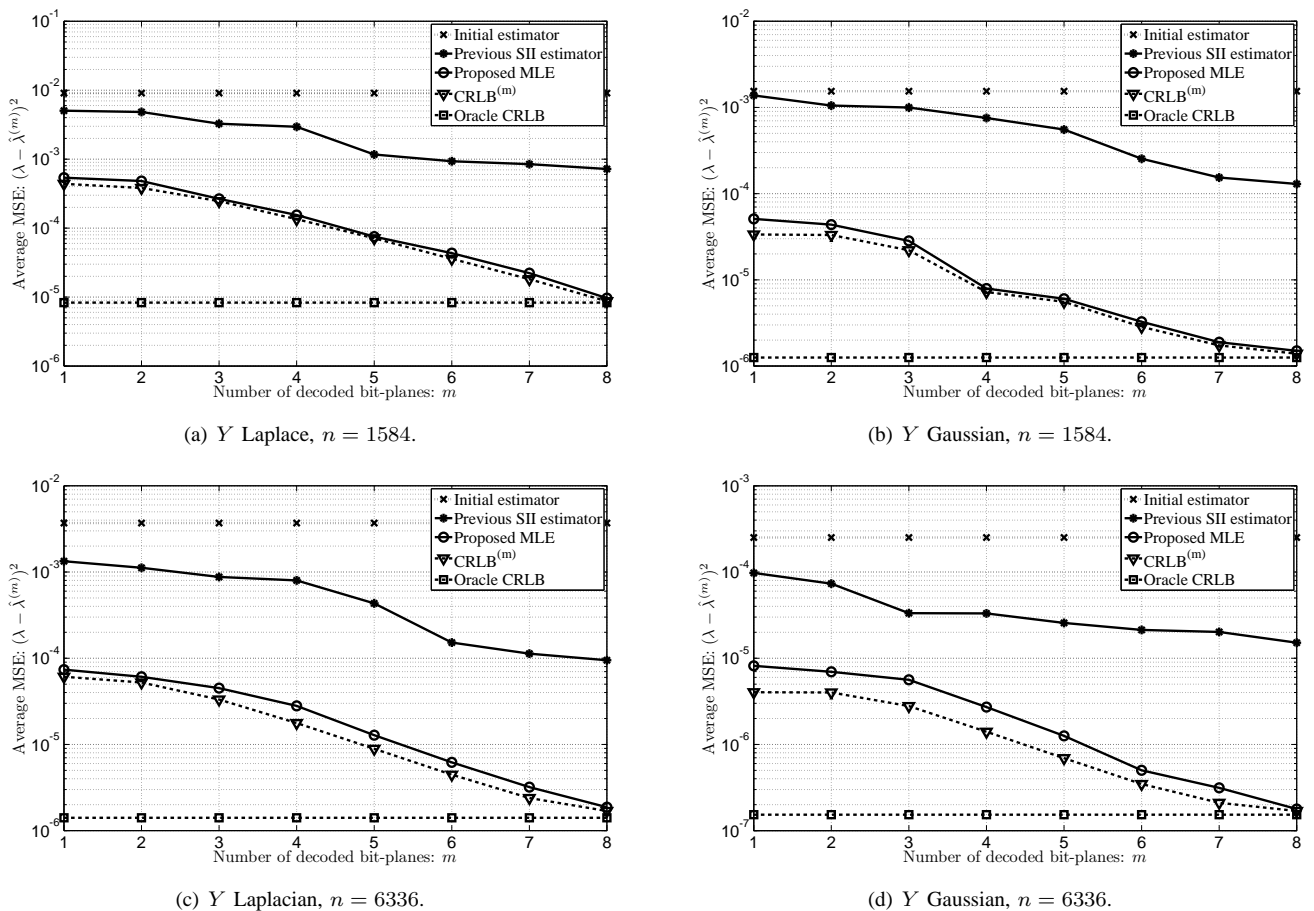


Fig. 2. The average MSE of the proposed MLE, our previous SII estimator from [33], and the initial estimator  $\hat{\lambda} = 1.5152$  with respect to the  $m$ -th stage CRLB $^{(m)}$  and the oracle CRLB, plotted versus the number of decoded bit-planes.

TABLE I  
BJØNTEGAARD RATE DELTAS,  $\Delta R(\%)$ , ON THE EDVC PERFORMANCE  
WHEN USING VARIOUS STATIONARITY BLOCK SIZES VERSUS THE  
PERFORMANCE OBTAINED ASSUMING NOISE STATIONARITY PER BAND  
PER FRAME.

$K$	1024	256	64	16	4
$(K_1 \times K_2)$	$(32 \times 32)$	$(16 \times 16)$	$(8 \times 8)$	$(4 \times 4)$	$(2 \times 2)$
Capsule Video 1	-0.18	-0.51	-0.99	-1.25	-0.99
Capsule Video 2	-0.03	-0.32	-0.60	-0.88	-0.76
Capsule Video 3	-0.18	-0.41	-0.64	-0.81	-0.48
Capsule Video 4	-0.11	-0.39	-0.59	-0.68	-0.18
Capsule Video 5	-0.20	-0.44	-0.62	-0.85	-0.39
$K$	384	64	24	16	4
$(K_1 \times K_2)$	$(24 \times 16)$	$(8 \times 8)$	$(6 \times 4)$	$(4 \times 4)$	$(2 \times 2)$
Endoscopy Video 1	-0.43	-0.95	-1.14	-1.21	-0.86
Endoscopy Video 2	-0.45	-0.73	-0.85	-0.94	-0.63
Endoscopy Video 3	-0.35	-0.77	-0.92	-1.13	-1.06

rate of 2Hz with a frame resolution of  $256 \times 256$  pixels. Contrary, each endoscopic sequence contains 100 frames with a resolution of  $480 \times 320$  pixels at a frame rate of 30Hz. The EDVC system was configured with a GOP of size two and all three Y, U and V components of the sequences were encoded, as detailed in [15].

Initially, we study the influence of the stationarity block size on the CCE accuracy and in turn, on the RD performance. The codec's configuration where a Laplace scaling

parameter  $\hat{\lambda}^{(m)}(\beta)$  is estimated per band of each frame (i.e.,  $n = K = 64 \times 64 = 4096$  for capsule endoscopic and  $n = K = 120 \times 80 = 9600$  for endoscopic sequences) is set as benchmark. Table I depicts the Bjøntegaard Deltas [43] (BD) on the performance of EDVC obtained with different stationarity block sizes,  $K_1 \times K_2$ , compared to the benchmark configuration. The results corroborate the trade-off between capturing the spatial variations of the noise and maintaining adequate support for the accurate derivation of the MLE. Overall, the best performance is achieved when regulating the stationarity block size to  $K = K_1 \times K_2 = 16$  coefficients.

The performance of the EDVC system with the proposed MLE method is thereafter compared with the performances obtained when using alternative up-to-date CCE techniques. The results in Fig. 4 and Table II show that the proposed MLE (using  $K = 16$ ) systematically outperforms the TRACE [25] method and our previous SID CCE [33] technique, with respective BD rate savings of up to 7.13% and 4.22%. TRACE has also been used as a benchmark to evaluate other CCE methods, including [27], [29], [33]. Unlike TRACE and our SID CCE method, the proposed MLE is asymptotically efficient per decoding stage. In addition, contrary to TRACE, the proposed method performs a bit-plane-per-bit-plane refinement of the CCE. This is a feature of other existing CCE methods as well, including [26], [27], [30], [31] and our SID CCE



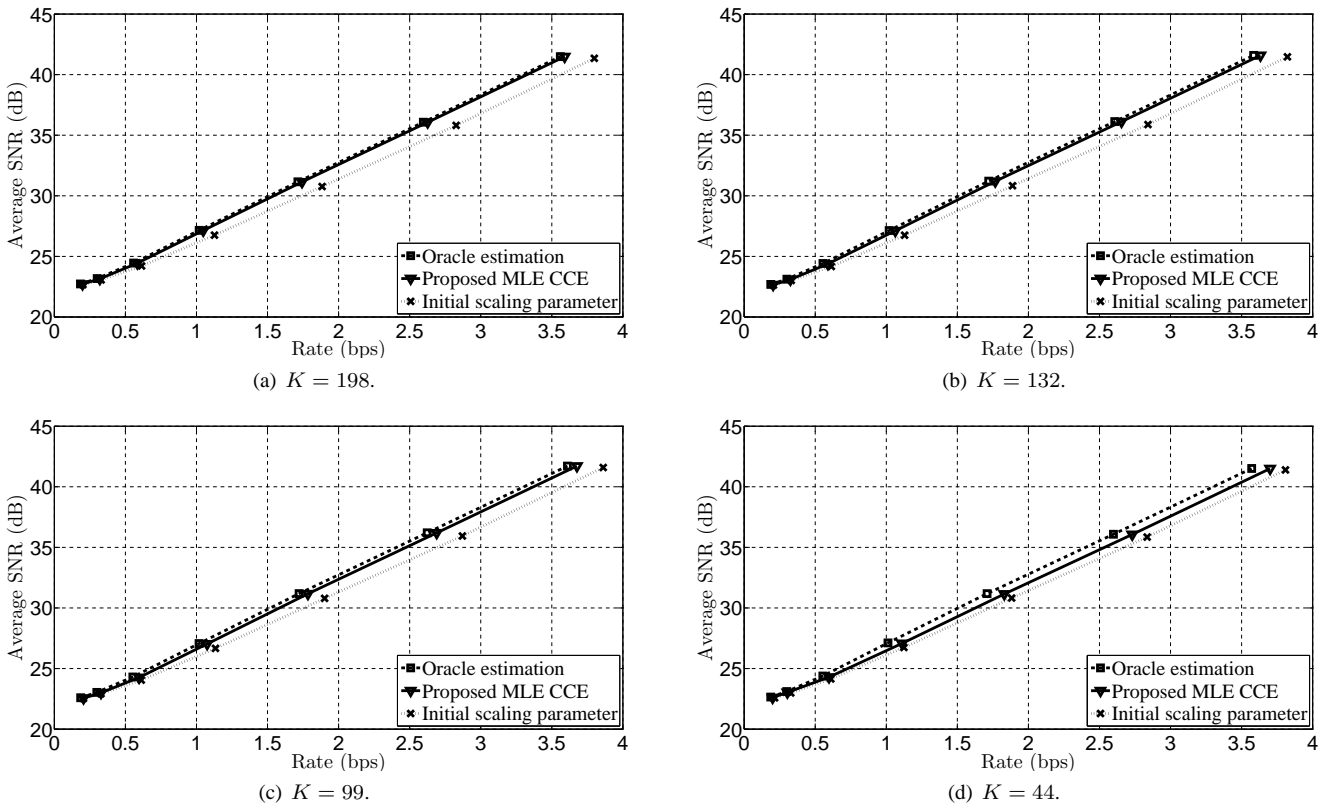


Fig. 3. RD performance of layered WZ coding with the proposed MLE. The side information is Laplacian, while various noise stationarity levels are tested.

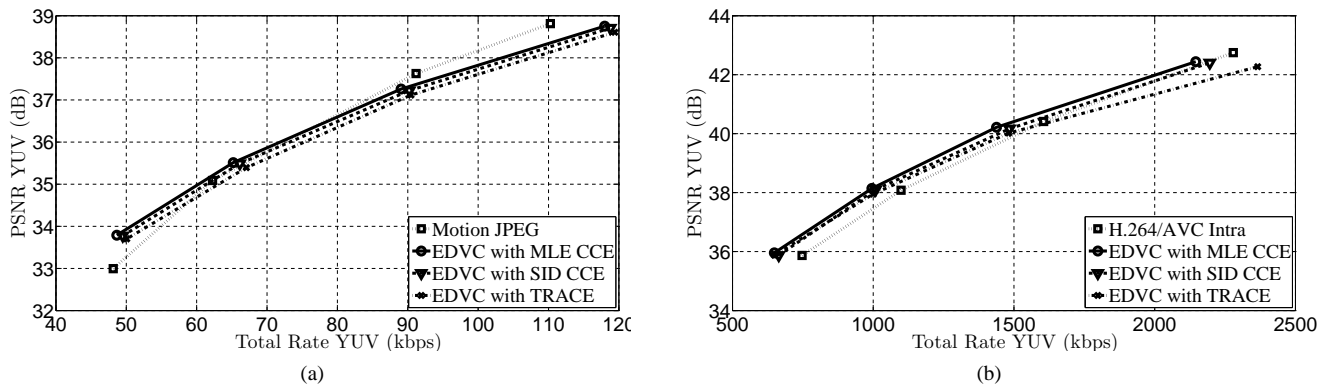


Fig. 4. Compression performance of EDVC using different online CCE methods versus Motion JPEG or H.264/AVC Intra: (a) Capsule Video 1; (b) Endoscopy Video 1. The rate is expressed in kilo-bits-per-second (kbps), while the distortion is quantified using the peak-signal-to-noise-ratio ( $\text{PSNR}_{\text{YUV}} = \frac{4 \times \text{PSNR}_Y + \text{PSNR}_U + \text{PSNR}_V}{6}$ ). The key and the hash frames of EDVC are encoded with (a) JPEG, or (b) H.264/AVC Intra.

technique. Compared with our SID CCE, however, the proposed MLE can adapt better to small noise stationarity levels, thereby improving the accuracy of the correlation estimation.

The results in Fig. 4(a) show that EDVC, equipped with the proposed MLE technique, yields comparable performance to Motion JPEG despite the highly-complex motion conditions encountered in capsule endoscopic video content. These conditions are due to the erratic motion of the capsule in the gastrointestinal tract and the low acquisition frame rates. Yet, when temporally fluent endoscopic video content is encoded [see Fig. 4(b)], EDVC introduces a notable BD rate gain of 8.80% over H.264/AVC Intra (Main Profile with CABAC).

### C. Complexity Assessment

The complexity of the proposed MLE technique is comparatively evaluated using execution time measurements obtained with the software implementation of our EDVC system<sup>7</sup>. This complexity examination methodology is commonly followed in the literature [25], [30], [33]. Table III depicts the execution time per WZ frame for the proposed MLE (configured with  $K = 4096$  or  $K = 16$ ) and our previous SID CCE [33] method. The total decoding time and the number of feedback

<sup>7</sup>Experiments were executed on an Intel® Core™ i7 CPU running at 2.20GHz with 16GB of RAM. Our EDVC was written in C++, compiled with Visual Studio 2008 and running in release mode under Windows 7.

TABLE II

BJØNTEGAARD DELTAS ON THE PERFORMANCE OF EDVC WHEN USING THE PROPOSED MLE METHOD VERSUS UP-TO-DATE CCE TECHNIQUES.

	vs. TRACE [25]		vs. SID CCE [33]	
	$\Delta R(\%)$	$\Delta \text{PSNR}(\text{dB})$	$\Delta R(\%)$	$\Delta \text{PSNR}(\text{dB})$
Capsule Video 1	-4.36	0.25	-1.82	0.11
Capsule Video 2	-5.15	0.33	-1.48	0.10
Capsule Video 3	-5.70	0.32	-1.89	0.11
Capsule Video 4	-5.04	0.29	-2.96	0.18
Capsule Video 5	-7.13	0.40	-4.22	0.23
Endoscopy Video 1	-4.67	0.24	-3.11	0.17
Endoscopy Video 2	-6.16	0.31	-3.50	0.21
Endoscopy Video 3	-5.05	0.25	-3.87	0.23

TABLE III

DECODING EXECUTION TIME (SEC) AND NUM. OF FEEDBACK REQUESTS PER WZ-FRAME OF THE EDVC SYSTEM CONFIGURED WITH DIFFERENT CCE TECHNIQUES; (A) CAPSULE VIDEO 4, (B) CAPSULE VIDEO 5.

	SID CCE [33]			MLE ( $K = 4096$ )			MLE ( $K = 16$ )		
	CCE time	Dec. time	#FB reqs.	CCE time	Dec. time	#FB reqs.	CCE time	Dec. time	#FB reqs.
(a) RD pt-1	0.03	27.40	66.5	0.09	25.86	54.60	0.11	25.29	54.37
RD pt-2	0.23	32.05	96.62	0.46	30.25	86.83	0.59	29.39	85.03
RD pt-3	0.42	50.43	148.17	0.64	46.72	139.52	0.90	44.03	136.06
RD pt-4	0.76	62.23	193.15	0.87	57.43	185.00	1.19	52.78	180.75
(b) RD pt-1	0.03	28.53	68.62	0.08	26.89	58.09	0.10	26.56	56.25
RD pt-2	0.18	33.74	98.52	0.22	31.40	93.37	0.37	30.77	91.38
RD pt-3	0.34	52.36	158.08	0.38	48.58	147.71	0.63	46.27	145.50
RD pt-4	0.72	63.80	204.08	0.79	58.93	196.15	1.10	55.02	193.25

requests per WZ frame of EDVC, configured with each of the assessed CCE techniques, is also provided.

The results indicate that the correlation estimation generally has a minor impact on the total decoding computational demands compared to other components of the decoder, particularly, the side information generation and the LDPCA decoding. Contrasting the complexities of the different CCE methods, it can be observed that, when the noise scaling parameter is estimated at a DCT band level (i.e.,  $K = 4096$ ), the proposed MLE induces a complexity comparable to that of our previous SID technique in [33]. Lowering the value of  $K$  increases the complexity of the proposed method, as the MLE algorithm is run per stationarity block of each band.

Nevertheless, incorporating the proposed MLE (especially using  $K = 16$ ) reduces significantly the overall decoding complexity of EDVC compared to employing our prior SID CCE method. In effect, as shown in Table III, by improving the CCE accuracy over our prior SID technique, the proposed MLE method induces fewer feedback channel requests and LDPCA decoding iterations. Hence, the feedback channel rate and the associated latency are also reduced.

## VII. CONCLUSION

A novel CCE technique tailored to generic layered WZ coding has been proposed in this paper. Unlike existing methods, e.g., [30], [32], [33], the proposed approach derives a novel MLE of the scaling parameter of the Laplacian correlation noise. Adhering to quality scalable WZ coding, the derived estimator is progressively refined per decoding stage. The MSE of the proposed estimator is proven to approximate the derived  $m$ -th stage CRLB<sup>(m)</sup>, converging to the oracle CRLB at high

decoding stages. It is also shown that the proposed MLE delivers an RD performance close to offline estimation. When incorporated into our EDVC system, the proposed method systematically yields superior RD performance over state-of-the-art techniques. Specifically, when endoscopic video content is encoded, the proposed method brings respective BD [43] rate savings of up to 7.13% and 4.22% over TRACE [25] and our previous SID CCE [33] technique. Finally, as verified by experimentation, the proposed algorithm is computationally efficient. In fact, on account of the improved accuracy in CCE, the LDPCA decoder converges faster, leading to reduced decoding delays that benefit applications like wireless capsule endoscopy.

## APPENDIX A

CALCULATION OF  $\frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda)$  AND  $\frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda)$

Applying the natural logarithm to the three cases of the pmf in (5) and then taking the first partial derivative with respect to  $\lambda$  yields the three forms of  $\frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda)$  shown in (11a), (11b) and (11c), where  $\Delta_i^{(m)}$  is the quantization interval of the  $i$ -th quantization cell,  $\theta_i^{(m)} = y_i - u_i^{(m)}$ , and  $\eta_i^{(m)} = v_i^{(m)} - y_i$ . Next, applying the partial derivative with respect to  $\lambda$  on (11a), (11b) and (11c) gives the three forms of  $\frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda)$  expressed by (12a), (12b) and (12c).

## APPENDIX B

EXISTENCE AND UNIQUENESS OF THE MLE

*Lemma 1:* The  $m$ -th stage conditional log-likelihood function  $\ln \mathcal{L}(\lambda)$  for a Laplace correlation channel is concave at  $\lambda \in (0, +\infty)$ .

*Proof:* Taking the second partial derivative of  $\ln \mathcal{L}(\lambda)$ , defined by (7), with respect to  $\lambda$ , yields

$$\begin{aligned} \frac{\partial^2}{\partial \lambda^2} \ln \mathcal{L}(\lambda) &= \frac{\partial^2}{\partial \lambda^2} \ln \prod_{i=1}^K p(q_i^{(m)}|y_i; \lambda) \\ &= \sum_{i=1}^K \frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda), \end{aligned} \quad (13)$$

where each term  $\frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda)$  is given by (11a), (11b) or (11c) depending on whether  $y_i < u_i^{(m)}$ ,  $y_i \geq v_i^{(m)}$  or  $u_i^{(m)} \leq y_i < v_i^{(m)}$ , respectively. Since  $\lambda \in (0, +\infty)$  and  $\Delta_i^{(m)} > 0$ , the summation terms in (13) that may take the form of (12a) or (12b) are evidently strictly negative. These terms correspond to  $(q_i^{(m)}, y_i)$  pairs for which  $y_i < u_i^{(m)}$  or  $y_i \geq v_i^{(m)}$ . We now turn to the terms that may follow the form of (12c). Since for these terms  $u_i^{(m)} \leq y_i < v_i^{(m)}$ , then by definition  $\theta_i^{(m)} \geq 0$  and  $\eta_i^{(m)} > 0$ , meaning that  $0 \leq e^{-\lambda \theta_i^{(m)}} \leq 1$  and  $0 \leq e^{-\lambda \eta_i^{(m)}} < 1$ . As a consequence, for the denominator of the first fraction in (12c) it holds that  $2 - e^{-\lambda \theta_i^{(m)}} - e^{-\lambda \eta_i^{(m)}} > 0$ . Namely, the terms in (13) that may follow (12c) are also strictly negative. As a consequence,  $\frac{\partial^2}{\partial \lambda^2} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) < 0$ . Therefore, according to the second derivative test,  $\ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$  is concave at  $\lambda \in (0, +\infty)$ . ■

*Proof of Theorem 1:* From Lemma 1, the function  $\frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$  is strictly monotonically decreasing at  $\lambda \in (0, +\infty)$ . Furthermore, the one-sided limit of the function at zero is

$$\lim_{\lambda \rightarrow 0^+} \frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = \sum_{i=1}^K \lim_{\lambda \rightarrow 0^+} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = +\infty, \quad (14)$$

since  $\lim_{\lambda \rightarrow 0^+} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = +\infty$  for all possible forms of  $\frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda)$ , given by (11a), (11b) or (11c).

Concerning the limit at infinity, we have: If  $y_i < u_i^{(m)}$ , then  $\lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = \theta_i^{(m)} < 0$ . Alternatively, if  $y_i > v_i^{(m)}$ , then  $\lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = \eta_i^{(m)} < 0$ . If  $u_i^{(m)} \leq y_i \leq v_i^{(m)}$ , then  $\lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = 0$ . Hence, if there exists at least one side information sample  $y_i$  in  $\mathbf{y}$  for which  $y_i < u_i^{(m)}$  or  $y_i > v_i^{(m)}$ , then

$$\lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = \sum_{i=1}^K \lim_{\lambda \rightarrow +\infty} \frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) < 0, \quad (15)$$

Combining (14) and (15) with the fact that  $\frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda)$  is continuous and strictly monotonically decreasing proves that the equation  $\frac{\partial}{\partial \lambda} \ln p(\mathbf{q}^{(m)}|\mathbf{y}; \lambda) = 0$  has a unique and finite solution at  $\lambda \in (0, +\infty)$ . ■

## APPENDIX C

## ANALYTICAL SOLUTIONS OF THE MLE

*Proof of Proposition 1:* Provided that  $y_i \notin [u_i^{(m)}, v_i^{(m)}], \forall i = \{1, 2, \dots, K\}$ , one can divide the pairs  $(q_i^{(m)}, y_i), i \in \{1, 2, \dots, K\}$ , in two disjoint sets  $S_A$  and  $S_B$ , corresponding to the cases  $y_i < u_i^{(m)}$  and  $y_i \geq v_i^{(m)}$ , respectively. Under this condition, (8) can be developed as

$$\sum_{i \in S_A} \left[ y_i - v_i^{(m)} + \frac{\Delta_i^{(m)}}{1 - e^{-\lambda \Delta_i^{(m)}}} \right] + \sum_{i \in S_B} \left[ v_i^{(m)} - y_i + \frac{\Delta_i^{(m)}}{e^{\lambda \Delta_i^{(m)}} - 1} \right] = 0, \quad (16)$$

Given that the quantizer  $Q_m$  has uniform intervals, namely,  $\Delta_i^{(m)} = \Delta^{(m)}, \forall i \in \{1, 2, \dots, K\}$ , then equation (16) becomes

$$\frac{|S_A| \Delta^{(m)}}{1 - e^{-\lambda \Delta^{(m)}}} + \frac{|S_B| \Delta^{(m)}}{e^{\lambda \Delta^{(m)}} - 1} = \sum_{i \in S_A} [v_i^{(m)} - y_i] + \sum_{i \in S_B} [y_i - v_i^{(m)}], \quad (17)$$

where  $|S_A|, |S_B|$  denote the cardinalities of the sets  $S_A$  and  $S_B$ . It holds that:  $|S_A| + |S_B| = K$ . Next, solving (17) for  $\lambda$  gives

$$\hat{\lambda}^{(m)} = \frac{1}{\Delta^{(m)}} \times \ln \frac{\sum_{i \in S_A} [v_i^{(m)} - y_i] + \sum_{i \in S_B} [y_i - v_i^{(m)}] + |S_B| \Delta^{(m)}}{\sum_{i \in S_A} [v_i^{(m)} - y_i] + \sum_{i \in S_B} [y_i - v_i^{(m)}] - |S_A| \Delta^{(m)}}. \quad (18)$$

$$\frac{\partial}{\partial \lambda} \ln p(q_i^{(m)}|y_i; \lambda) = \begin{cases} y_i - v_i^{(m)} + \frac{\Delta_i^{(m)}}{1 - e^{-\lambda \Delta_i^{(m)}}}, & y_i < u_i^{(m)}, & (11a) \\ v_i^{(m)} - y_i + \frac{\Delta_i^{(m)}}{e^{\lambda \Delta_i^{(m)}} - 1}, & y_i \geq v_i^{(m)}, & (11b) \\ \frac{\theta_i^{(m)} e^{-\lambda \theta_i^{(m)}} + \eta_i^{(m)} e^{-\lambda \eta_i^{(m)}}}{2 - e^{-\lambda \theta_i^{(m)}} - e^{-\lambda \eta_i^{(m)}}}, & u_i^{(m)} \leq y_i < v_i^{(m)}. & (11c) \end{cases}$$

$$\frac{\partial^2}{\partial \lambda^2} \ln p(q_i^{(m)}|y_i; \lambda) = \begin{cases} -\frac{[\Delta_i^{(m)}]^2 e^{-\lambda \Delta_i^{(m)}}}{[1 - e^{-\lambda \Delta_i^{(m)}}]^2}, & y_i < u_i^{(m)}, & (12a) \\ -\frac{[\Delta_i^{(m)}]^2 e^{\lambda \Delta_i^{(m)}}}{[e^{\lambda \Delta_i^{(m)}} - 1]^2}, & y_i \geq v_i^{(m)}, & (12b) \\ -\frac{[\theta_i^{(m)}]^2 e^{-\lambda \theta_i^{(m)}} + [\eta_i^{(m)}]^2 e^{-\lambda \eta_i^{(m)}}}{2 - e^{-\lambda \theta_i^{(m)}} - e^{-\lambda \eta_i^{(m)}}} - \left[ \frac{\theta_i^{(m)} e^{-\lambda \theta_i^{(m)}} + \eta_i^{(m)} e^{-\lambda \eta_i^{(m)}}}{2 - e^{-\lambda \theta_i^{(m)}} - e^{-\lambda \eta_i^{(m)}}} \right]^2, & u_i^{(m)} \leq y_i < v_i^{(m)}. & (12c) \end{cases}$$

Alternatively, if  $|S_B| = 0$  or  $|S_A| = 0$ , namely, if  $y_i < u_i^{(m)}$  or  $y_i \geq v_i^{(m)}$ ,  $\forall i \in \{1, 2, \dots, K\}$  then

$$\hat{\lambda}^{(m)} = \frac{1}{\Delta^{(m)}} \ln \left[ 1 + \frac{K\Delta^{(m)}}{\sum_{i=1}^K [y_i - v_i^{(m)}]} \right]^{\mp 1}, \quad (19)$$

where the  $-$  and  $+$  on the exponent corresponds to the first and second case, respectively. ■

**CRLB<sup>(m)</sup> for the Analytical Forms:** Provided that the conditions listed in Proposition 1 hold, then replacing the second derivative forms of (12a) and (12b) in (10) gives

$$\text{CRLB}^{(m)} = \left[ \frac{|S_A|\Delta^{(m)}e^{-\lambda\Delta^{(m)}}}{(1 - e^{-\lambda\Delta^{(m)}})^2} + \frac{|S_B|\Delta^{(m)}e^{\lambda\Delta^{(m)}}}{(e^{\lambda\Delta^{(m)}} - 1)^2} \right]^{-1}. \quad (20)$$

Setting  $|S_A| = K$  and  $|S_B| = 0$  or  $|S_B| = K$  and  $|S_A| = 0$  in (20) yields the respective CRLB<sup>(m)</sup> for the case  $y_i < u_i^{(m)}$  or  $y_i \geq v_i^{(m)}$ ,  $\forall i \in \{1, 2, \dots, K\}$ . ■

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